

# LANCOM™ Techpaper

## Routing Performance

### Introduction

Applications for communications and entertainment are increasingly based on IP networks. In order to ensure that the necessary bandwidth performance can be provided reliably, it is important for the infrastructure's networking components to be tested thoroughly and intensively. In this Techpaper, LANCOM Systems presents the methods of measuring the performance of routing and VPN systems for central sites and VPN gateways.

We examine a variety of aspects for consideration when measuring the performance of a router. This includes the transmission speeds of connections between the LAN and the Internet (WAN), and the internal data transmission in the network (LAN-LAN). Many business processes rely on secure WAN connections, which is why we are focusing on determining the performance of encrypted data connections over VPN.

### Test system

All of the performance values were measured in the LANCOM test laboratory. Tests were conducted with an IXIA test system. IXIA uses so-called test suites, which enable the simulation of different applications. This allows, for example, the investigation of data throughput over automatically established VPN tunnels, or the measurement of pure LAN-WAN routing performance for unidirectional and bidirectional data connections. IXIA is a leading supplier of systems which test IP-based services and infrastructures. Test systems from IXIA are employed all over the world by network component manufacturers and other organizations to help assure the functionality and reliability of complex IP networks, devices, and applications. The measurement of data transmission itself uses either a fixed frame size or a combination of frame sizes which reflects a typical flow of data. These combinations are known as "Internet Mix", or IMIX for short. The type of

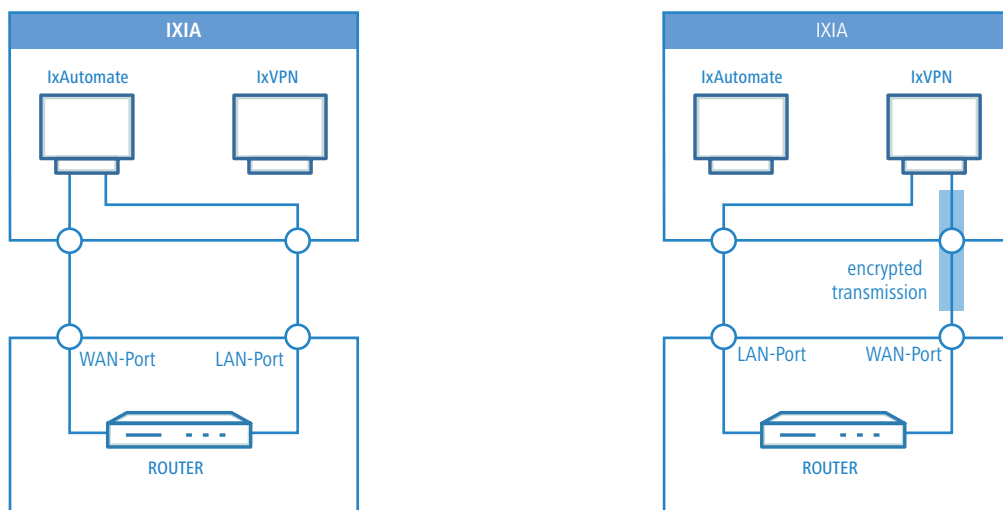


Figure 1: IXIA Test System for routing connections and encrypted VPN connections between LAN and WAN

IMIX which is applied significantly affects the test results because packet size has a strong influence on a connection's performance. By selecting the appropriate ports on the router being tested, it is possible to test connections between the LAN and the WAN, and also purely LAN-LAN connections.

### Routing Performance (UDP)

The measurement of routing performance involves the determination of the maximum data throughput which can be achieved before a router starts rejecting packets. Measurement uses UDP packets of various sizes in order to simulate the performance with different applications. Ethernet frame sizes range from 64 bytes for the smallest

Device	64	128	256	512	1024	1280	1518	frame rate
LANCOM 1781EW	5.960	12.256	27.470	47.949	91.971	114.922	134.957	11546.660
LANCOM 1781EF+	48.411	95.952	183.119	347.826	674.572	863.698	983.161	87176.857
LANCOM 7100+ VPN	62.622	125.000	237.037	467.153	909.414	984.615	986.996	108944.698
LANCOM 9100+ VPN	87.912	175.342	343.164	680.851	980.843	984.615	986.996	139125.706

Table 1: LAN-LAN routing - throughput [Mbps] per frame size [byte] and average frame rate [fps]

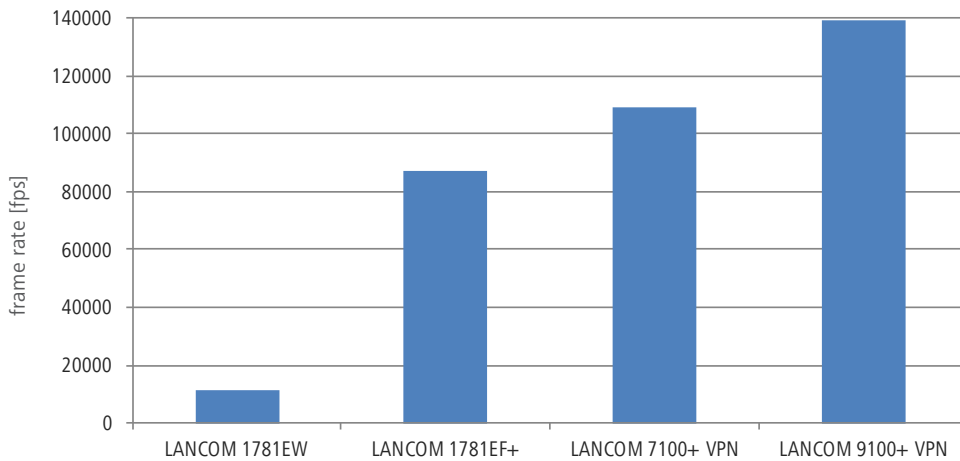


Diagram 1: LAN-LAN routing - frame rate

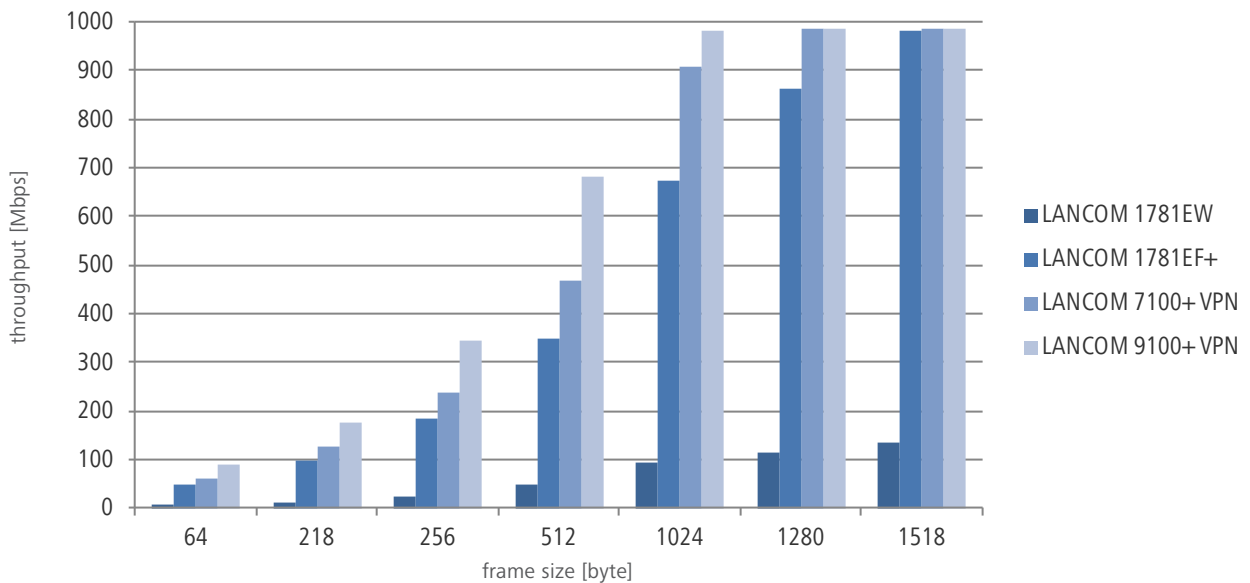


Diagram 2: LAN-LAN routing - throughput

to 1518 bytes for the largest frames. Tests on different router models demonstrate the influence of the different hardware platforms (processor, interfaces).

Measurements initially determine the frame rate, which is a good performance indicator of the tested hardware. With normal routing, the frame rate is fairly constant even with different frame sizes. This is because only the header is inspected during routing, a process which is largely independent of the size of the frames being routed. For this reason, only the average frame rates are given in the tables.

The throughput for a certain frame size (or even a mix of sizes, see IMIX on page 6) can be approximately calculated by multiplication with the frame rate. When the frame rate is constant, data throughput depends directly on the frame size because the larger the frames, the larger is the data volume that can be transmitted. The maximum number of frames transmitted per second is limited by the performance of the interfaces and the transmission medi-

um. Measurement of the routing performance relates to the size of the Ethernet frames. To compare packet sizes for particular applications, it is necessary to subtract the header. For a frame of 512 bytes, the result is a UDP datagram size of 470 bytes (512 bytes - 18 bytes Ethernet header - 4 bytes FCS trailer - 20 bytes IP header) and, after subtracting the UDP header (8 bytes), the UDP payload is 462 bytes.

To investigate routing performance, in this paper two different applications are considered.

- For WAN-LAN routing, data received from the WAN is forwarded to a peer in the LAN.
- For LAN-LAN routing, data remains within the local-area network and is passed from one LAN port to another.

The measurements show that the throughput increases almost linearly with the frame size until the limit of the gigabit interface is reached.

Device	64	128	256	512	1024	1280	1518	frame rate
LANCOM 1781EW	5.216	10.144	20.844	41.837	83.347	103.862	122.419	10126.166
LANCOM 1781EF+	32.389	65.911	125.490	253.968	482.564	642.570	736.893	61891.646
LANCOM 7100+ VPN	45.455	90.909	182.077	360.056	721.127	882.759	985.714	87110.110
LANCOM 9100+ VPN	62.136	131.417	265.010	520.325	978.967	984.615	986.996	114722.741

Table 2: WAN-LAN routing - throughput [Mbps] per frame size [byte] and average frame rate [fps]

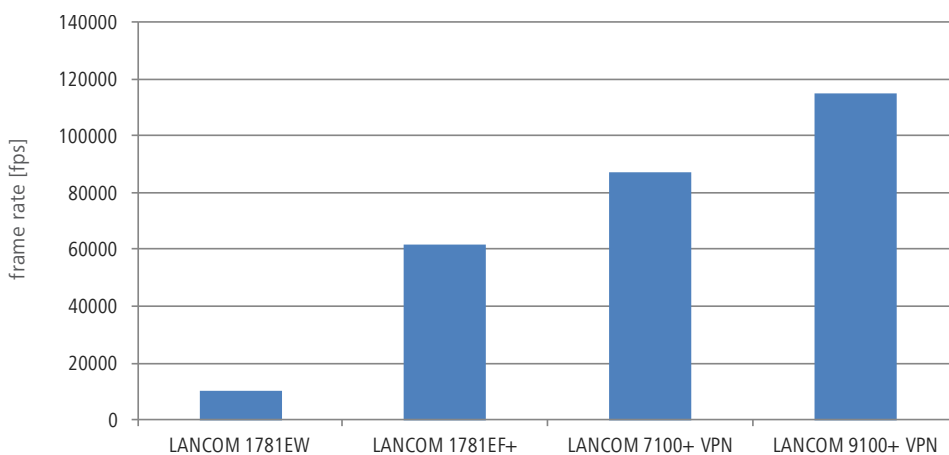


Diagram 3: WAN-LAN routing - frame rate

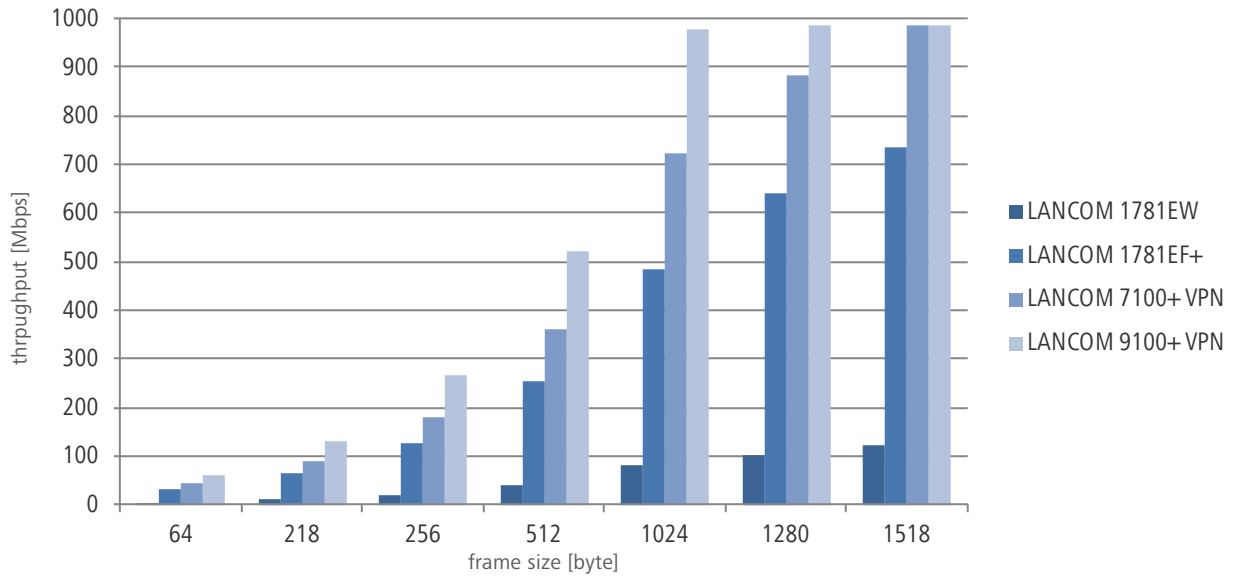


Diagram 4: WAN-LAN routing - throughput

### Routing Performance (TCP)

Using UDP, the maximum performance is easily measured. However, a large portion of data transfers utilizes TCP. Therefore, it is important to take a closer look at those scenarios as well.

The TCP measurement is accomplished by using iperf, a program used to measure the data throughput between two computers. They are connected via one LANCOM router. One computer acts as a server in the WAN, sending data packets to the other, which acts as a client in the LAN, representing a typical download scenario.

The TCP measurement done with iperf did not use NAT since no direct measurement with NAT is possible using iperf. The correlation between the test results and the results seen at a WAN connection including NAT in a live environment is very high. The causes are the differences of the traffic structure of iperf compared to an HTTP download and the LCOS 8.80 firewall which offers a similar performance with NAT or without.

Device	Throughput with 5 TCP sessions
LANCOM 1781EW	95.6
LANCOM 1781EF+	508.4
LANCOM 7100+ VPN	691.3
LANCOM 9100+ VPN	939.0

Table 3: WAN-LAN routing - throughput [Mbps]

The measurement was run using iperf 2.05. The TCP window size was set to 256 kb and five sessions were run simultaneously.

Both computers used in the test system have a similar hardware and software:

- Intel Core i7 cpu
- Intel PRO/1000 nic
- Ubuntu 11.04 / Kernel 2.6.38

The measurement was run using iperf 2.05. The TCP window size was set to 256 kb and five sessions were run simultaneously.

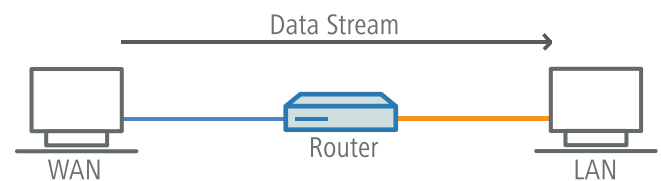


Figure 2: Schematic view of the test system

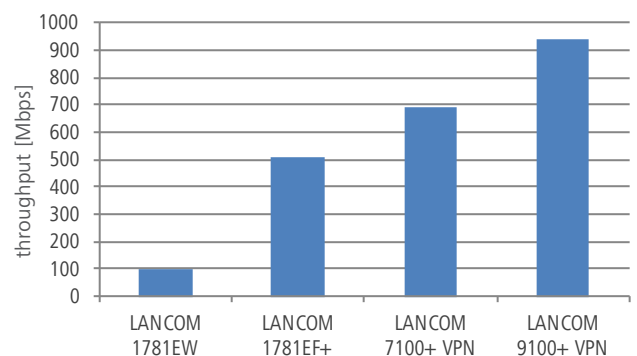


Diagram 5: TCP routing - throughput

### IPSec Routing Performance

Other than with pure routing performance, VPN and IPSec routing actually changes the frames which are being passed from one interface to the next. When data is encrypted for the VPN tunnel, the original frame is encapsulated and it is supplemented with additional information. This has two important effects when considering the performance of IPSec routing:

- Encrypted frames are larger than unencrypted frames. Consequently, any results have to indicate which

frame size was observed at which interface, and/or whether the frames were encrypted or unencrypted. The values presented here always relate to an unencrypted frame size. An IP packet of 42 bytes is transported unencrypted, e.g. in a frame of 64 bytes. In the event of AES encryption, the frame grows for example to 126 bytes (46 byte IP packet + 18 byte Ethernet + 4 byte FCS + 20 byte IP + 8 byte ESP + 16 byte initialization vector (IV) + 1 byte padding + 1 byte padding length + 12 byte authentication).

Device	64	128	256	512	1024	1280	1418	frame rate
LANCOM 1781EW	3.68	8.42	16.90	34.25	65.23	80.70	88.7	7955.471
LANCOM 1781EF+	18.70	38.05	76.99	144.42	277.00	339.58	371.76	35181.171
LANCOM 7100+ VPN	28.81	56.17	106.68	205.65	386.89	469.08	515.09	49980.280
LANCOM 9100+ VPN	39.96	76.63	149.19	286.13	530.13	645.38	712.00	69440.562

Table 4: IPSec routing (decryption) - throughput [Mbps] per frame size [byte] and average frame rate [fps]

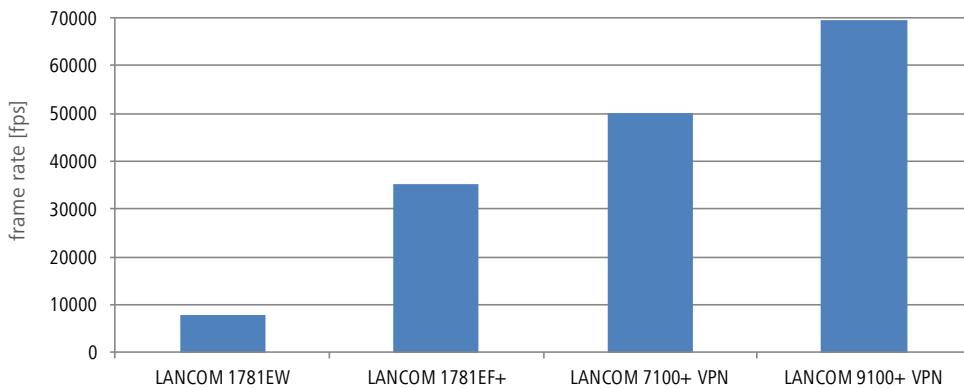


Diagram 6: IPSec routing (decryption) - frame rate

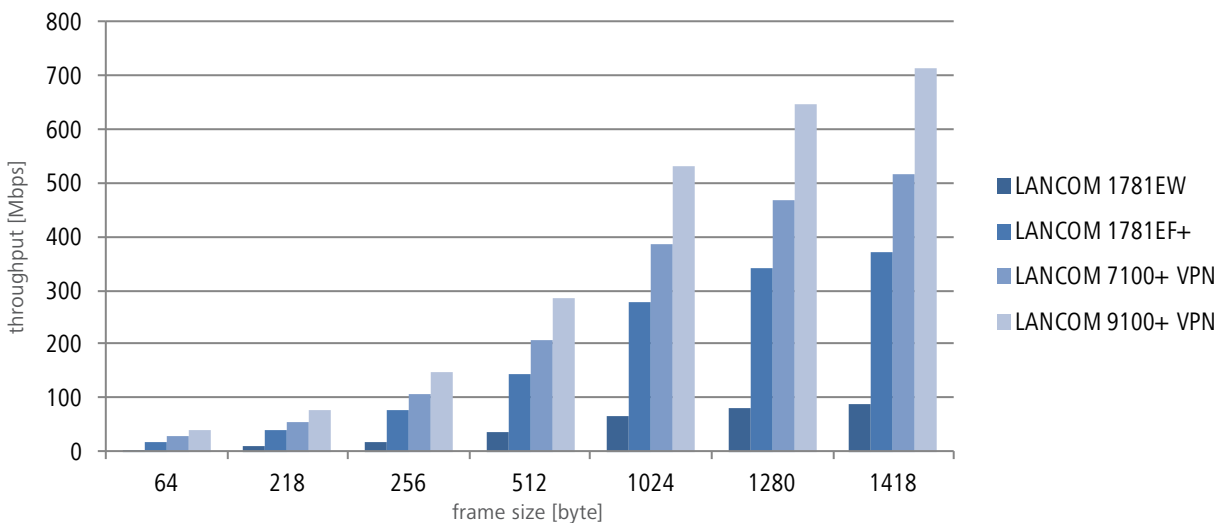


Diagram 7: IPSec routing (decryption) - throughput

■ The processes of encryption and decryption in the router take up computing time. These processes take place in two steps which, in the case of encryption, must be sequential. With decryption, on the other hand, these steps can be executed in parallel. Router models with VPN hardware acceleration provide significantly better performance with decryption than with encryption. This explains why the results display a significant difference in performance between the decryption and encryption directions.

All of the IPSec-routing values given here are for a single VPN tunnel. With up to 1000 tunnels established under laboratory conditions, the frame rate remained almost constant over all of the active tunnels. However, under actual operating conditions an increasing number of tunnels will cause the frame rate to drop due to the processes running for each tunnel (for example renewal of the key being used).

Device	64	128	256	512	1024	1280	1418	frame rate
LANCOM 1781EW	3.09	6.29	11.56	24.14	47.07	58.83	65.50	5854.40
LANCOM 1781EF+	14.72	30.38	60.65	119.17	232.60	283.89	316.57	28734.53
LANCOM 7100+ VPN	18.40	36.41	73.95	144.05	281.80	346.73	375.86	34739.80
LANCOM 9100+ VPN	26.10	52.34	103.42	203.68	390.31	472.64	519.23	48841.61

Table 5: IPSec routing (encryption) - throughput [Mbps] per frame size [byte] and average frame rate [fps]

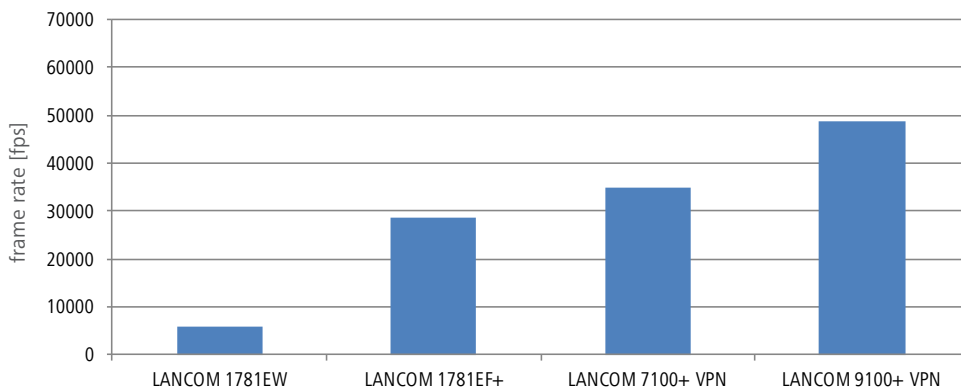


Diagram 8: IPSec routing (encryption) - frame rate

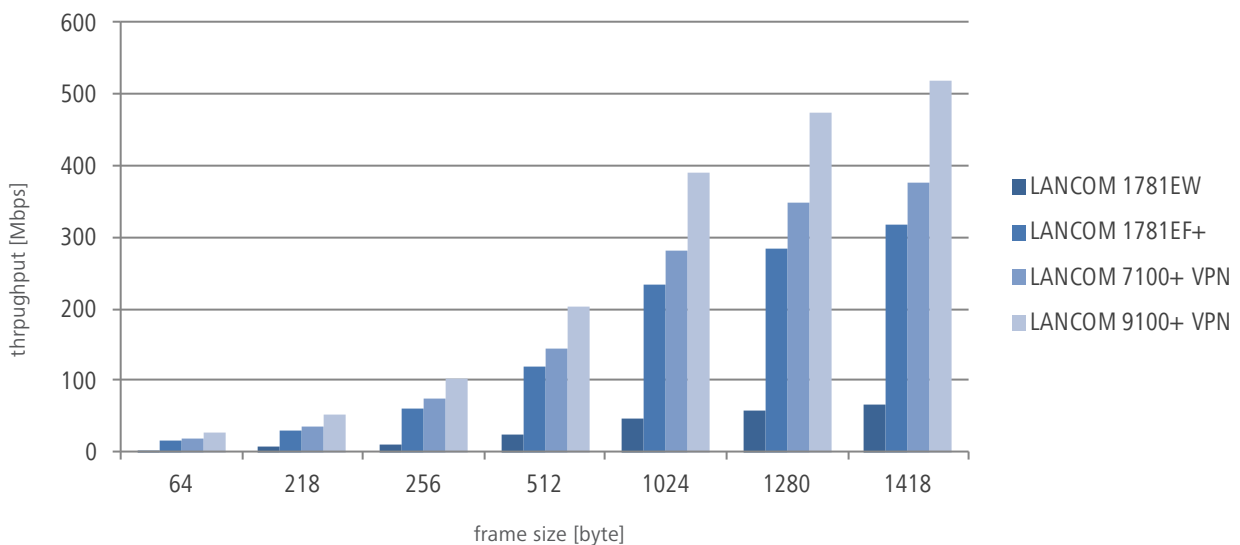


Diagram 9: IPSec routing (encryption) - throughput

Device	IMIX 0	IMIX 1	IMIX 2
LANCOM 1781EW	25.38	19.52	21.43
LANCOM 1781EF+	118.09	95.70	115.16
LANCOM 7100+ VPN	166.88	131.73	162.02
LANCOM 9100+ VPN	234.21	184.43	225.46

Table 6: IPSec routing (decryption) - throughput [Mbps]

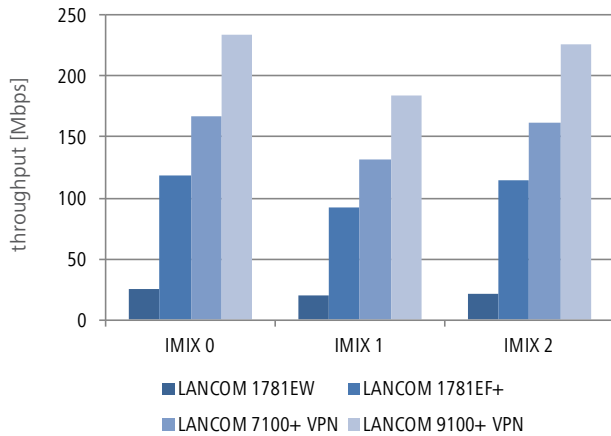


Diagram 10: IPSec routing (decryption)

Device	IMIX 0	IMIX 1	IMIX 2
LANCOM 1781EW	22.19	17.20	21.43
LANCOM 1781EF+	112.08	90.55	109.39
LANCOM 7100+ VPN	138.73	111.20	135.09
LANCOM 9100+ VPN	193.11	158.21	189.44

Table 6: IPSec routing (encryption) - throughput [Mbps]

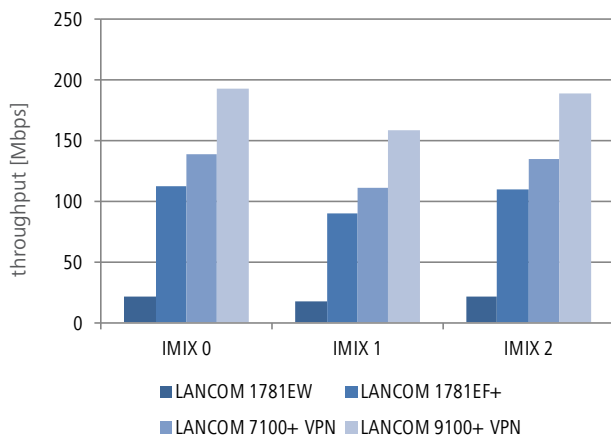


Diagram 11: IPSec routing (encryption)

### IPSec routing with different IMIXs (decryption and encryption)

As an alternative to measurement with fixed frame sizes, measurements were carried out with varying IMIX patterns. The IMIX patterns simulate “true” data traffic, which contains varying frame sizes. There are no fixed guidelines for the composition of the frame sizes, and for this reason the measurements supplemented the preset values in the IXIA test system (IMIX 0) with two other common patterns (IMIX 1 and IMIX 2). The patterns use the following frame compositions:

- IMIX 0: 45% 64 byte, 20% 128 byte, 5% 256 byte, 3% 512 byte, 2% 1024 byte, 1% 1280 byte, 24% 1364 byte.
- IMIX 1: 7x 64 byte, 4x 570 byte, 1x 1418 byte.
- IMIX 2: 58% 90 byte, 2% 92 byte, 24% 594 byte, 16% 1418 byte.

The maximum encrypted frame size transmittable via Ethernet is 1418 bytes (1518 bytes is the maximum IEEE802.3 frame size).

AES128-SHA encryption was used for the measurements, and the tunnels were established in the LAN-WAN direction. These measurements again show that data decryption is faster than encryption.